

# **Data Structures and Algorithms**

## **- Dynamic Hashing -**

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# Dynamic External Hashing

## ➤ Extendible Hashing

- ▶ use an access structure in addition to the file
- ▶ based on the result of the hash function to the search field
- ▶ similar to index but based on the search field.

## ➤ Linear Hashing

- ▶ do not need any access structure
- ▶ based on a sequence of hash functions

# Extendible Hashing(1)

- A directory can be stored on disk, and it expands or shrinks dynamically. Directory entries point to the disk blocks that contain the stored records.
- A directory of  $2^d$  bucket addresses, where  $d$  is called the **global depth** of the directory.
- The first  $d$  bits of a hash value as an index into the directory.
- Several directory locations with the same first  $d'$  (**local depth**) bit for their hash values may contain the same bucket address if all the records that hash to these locations fit in a single bucket.

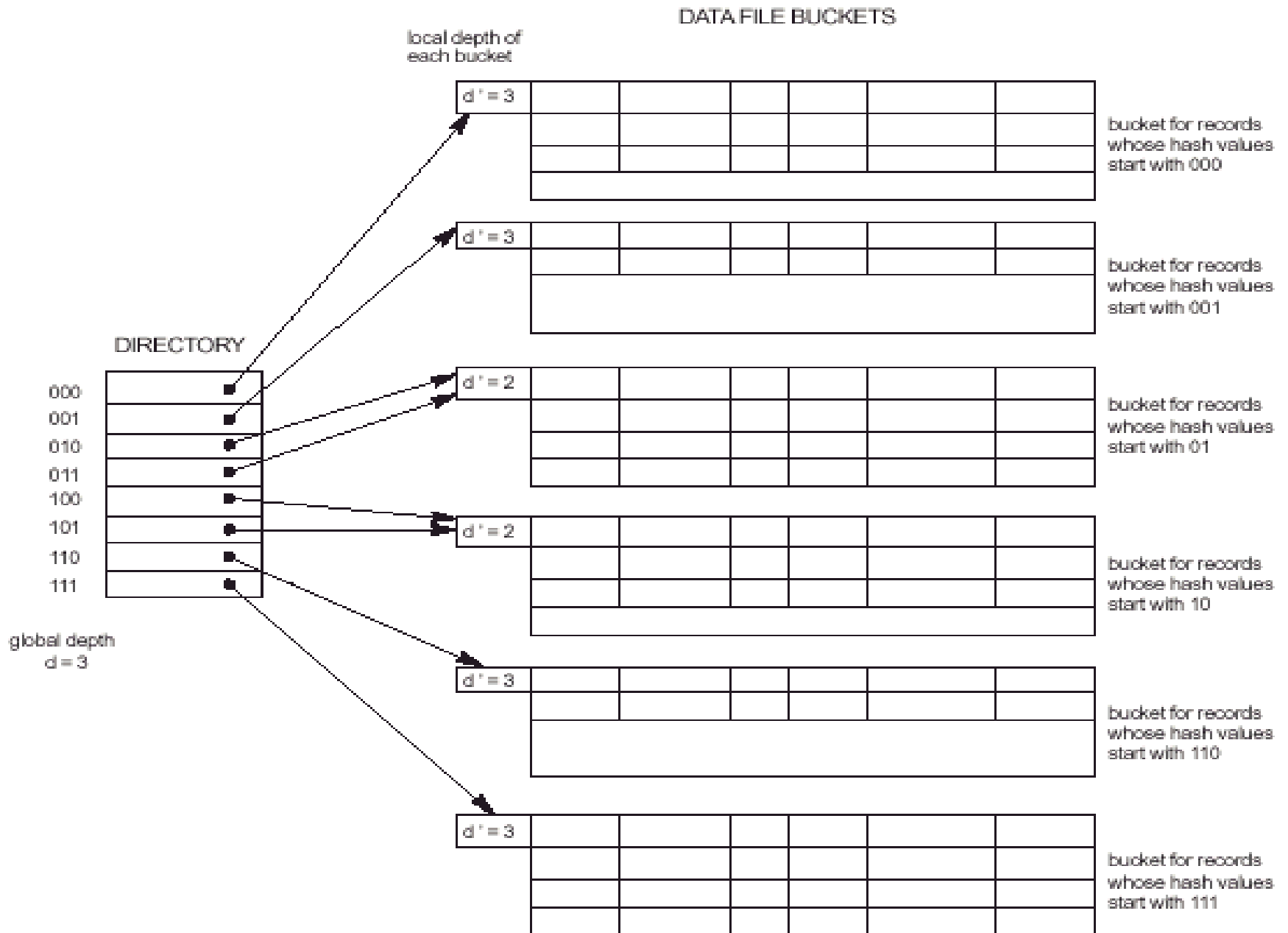


Fig 5.13 Structure of the extendible hashing scheme.

# Extendible Hashing(2)

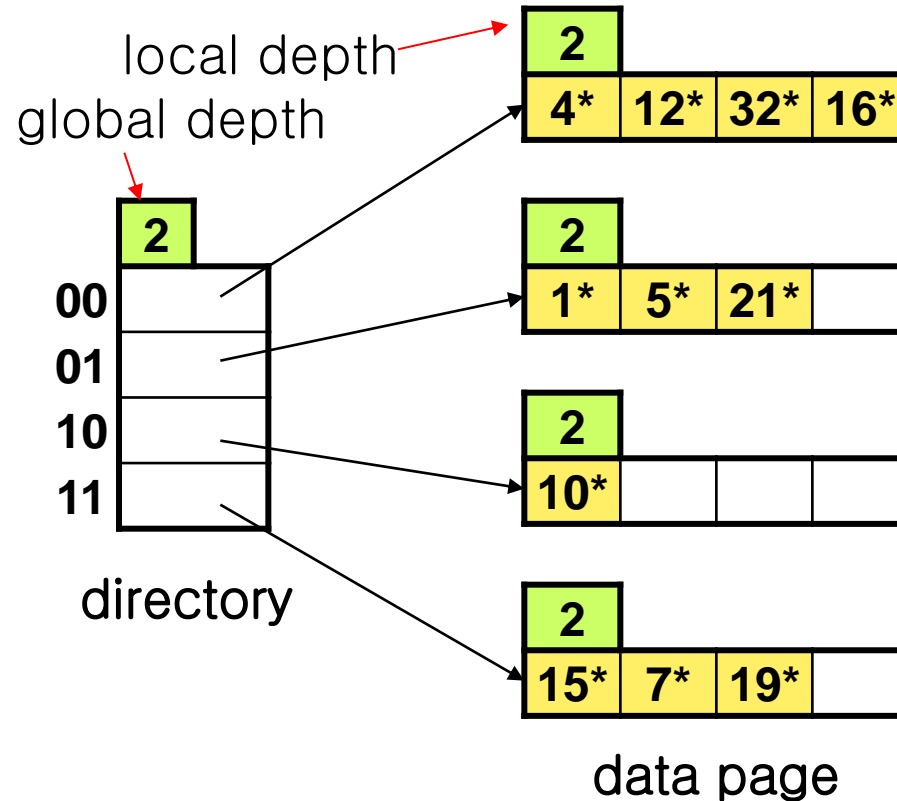
- Incrementing  $d$  by one
  - ▶ Doubling the number of entries in the directory
  - ▶ When a bucket, whose  $d'$  is equal to  $d$ , overflows.
- Decrementing  $d$  by one
  - ▶ Halving the number of entries in the directory
  - ▶ When  $d > d'$  for all buckets after some deletions.
- Does not require an overflow area.
- Two block accesses for record retrieval.
  - ▶ One for directory
  - ▶ One for bucket

# Managing directory

- When a bucket with  $d' = d$  overflows,
  - ▶ split the bucket
  - ▶ distribute records based on  $(d+1)$ th bit
  - ▶ double the directory
  - ▶ adjust the directory entries
  
- When  $d > d'$  for all the buckets,
  - ▶ halve the directory
  - ▶ adjust the directory entries

# Extendible hashing : 예제(1)

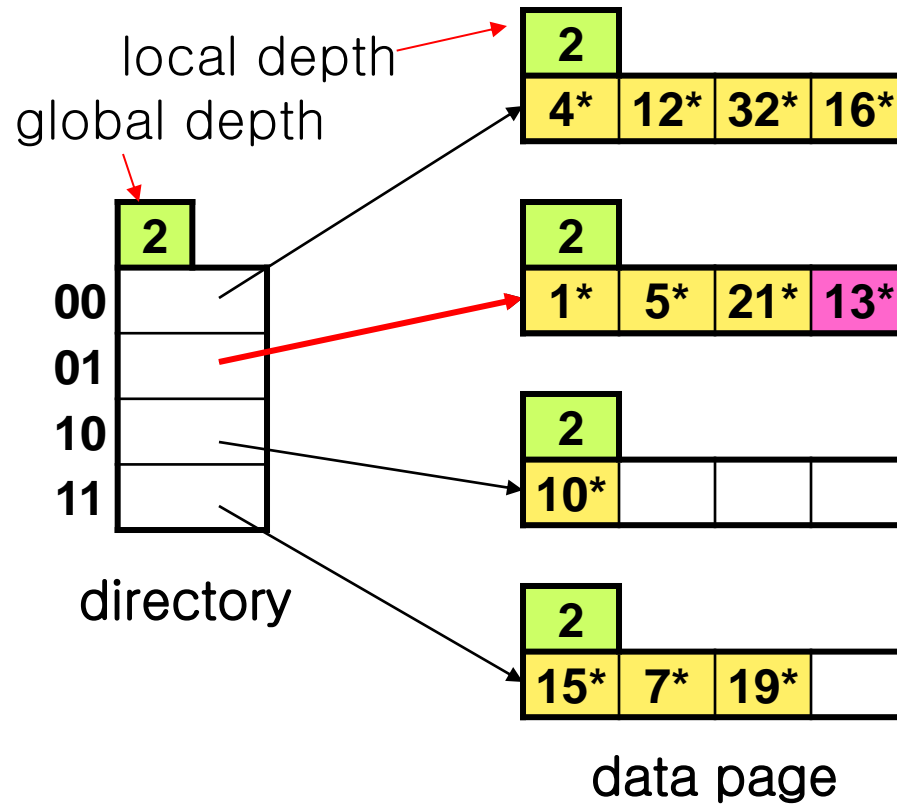
- Directory는 크기 4의 배열 (내용은 bucket에 대한 pointer)
- 각 bucket은 최대 4개의 data entry 저장
- $h(K)$  값의 binary 수의 **마지막 2 bit**를 directory에 적용



(예) insert 13  
 $13 \bmod 4 = 1(01)$

# Extendible hashing : 예제(2)

(예)  
insert 13 (01)

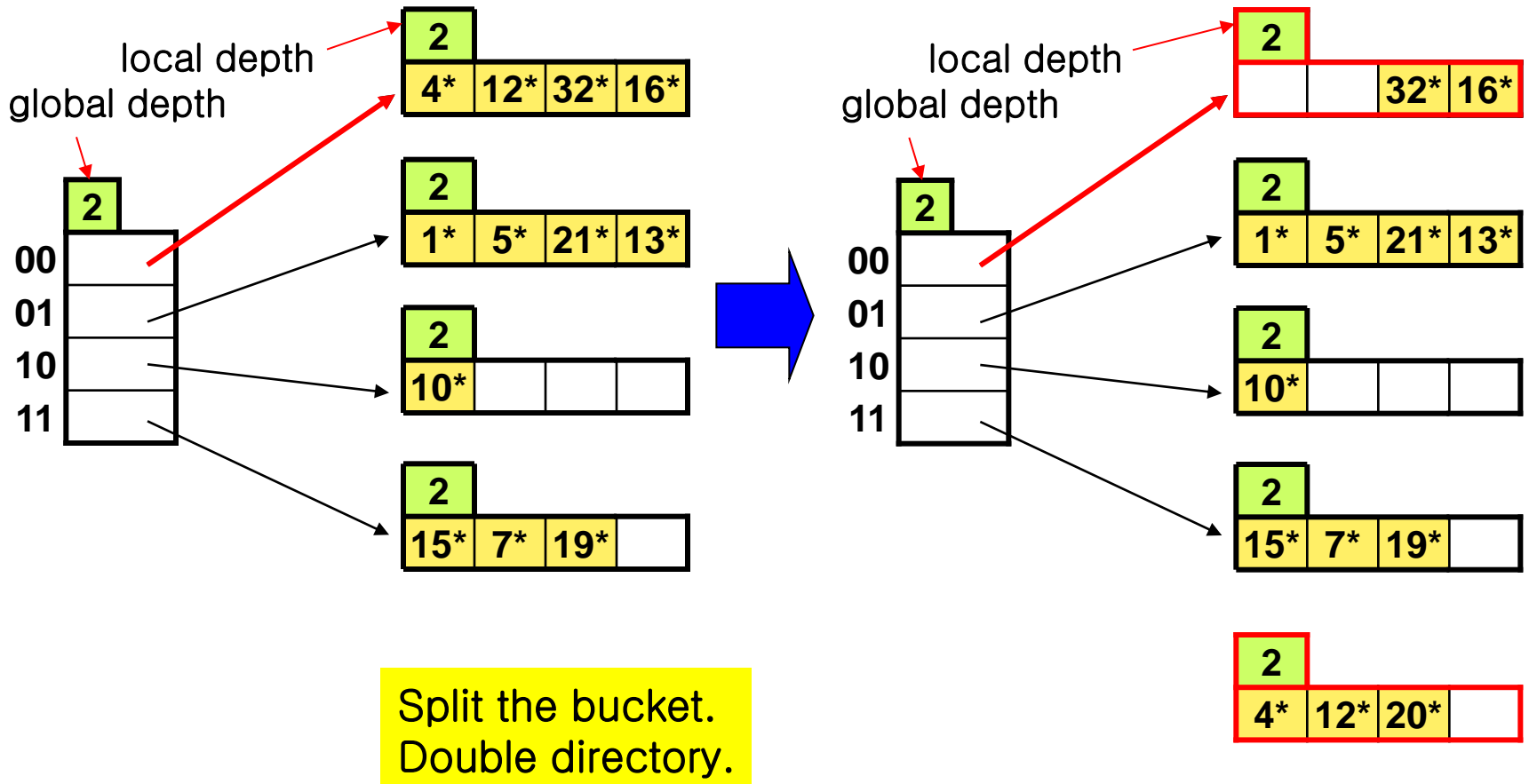




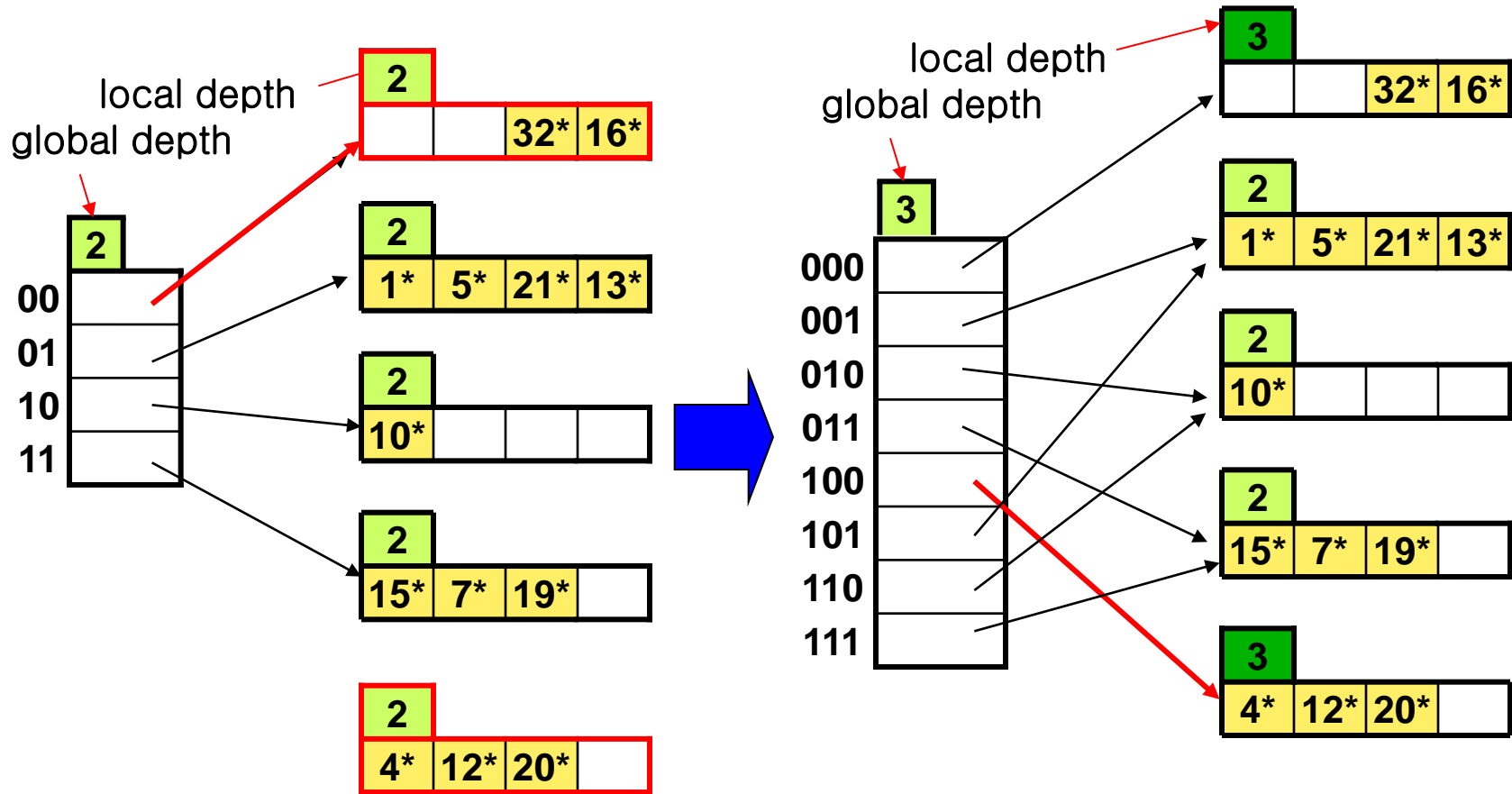
# Extendible hashing : 예제(3)

➤ full bucket에 data entry insert

(예) insert 20 ( $20 \bmod 4 = 0$  (00))

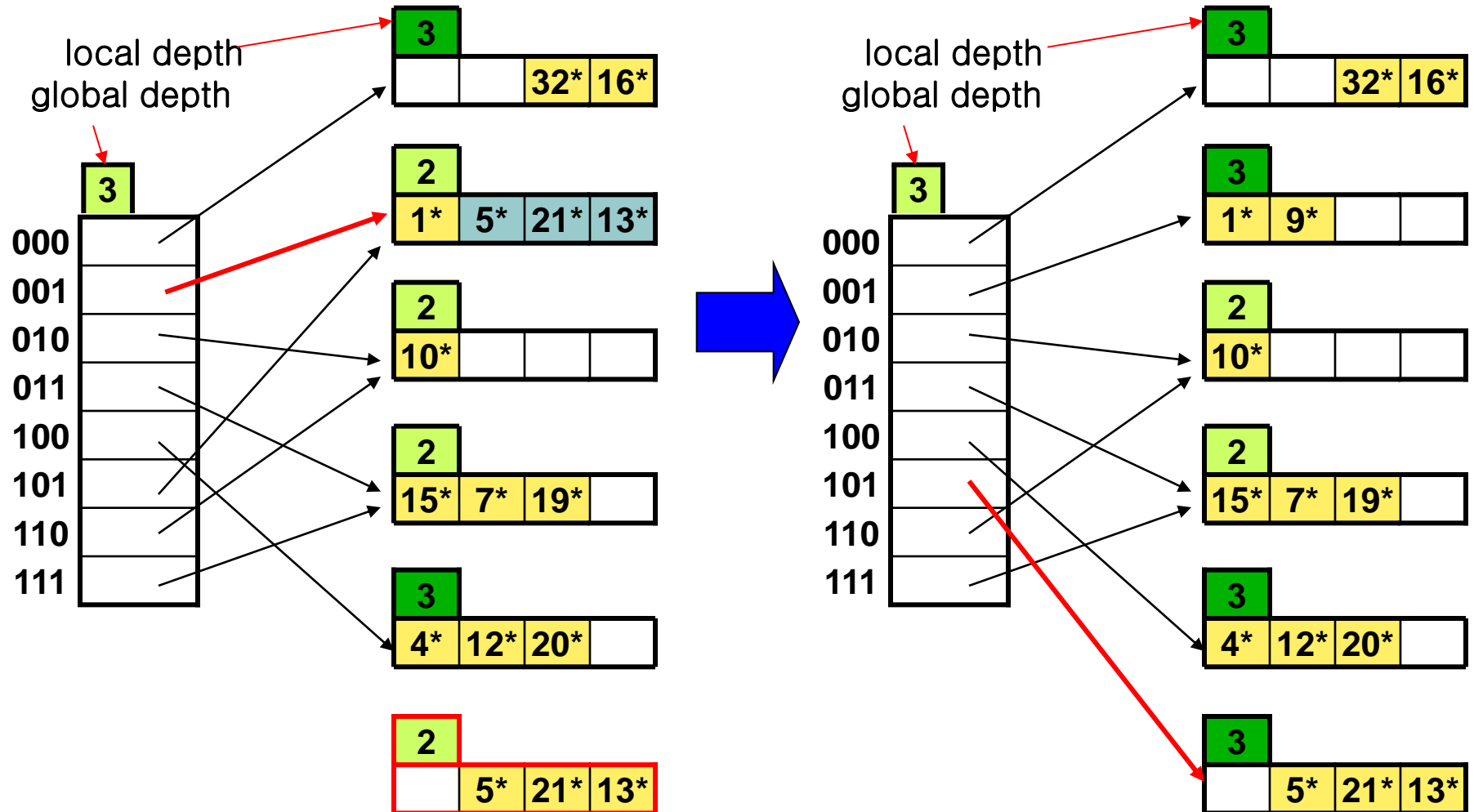


# Extendible hashing : 예제(4)



# Extendible hashing : 예제(5)

(예) insert 9 ( $9 \bmod 4 = 1$  (01))



# Properties

## ➤ Advantages

- ▶ Performance doesn't degrade as the file grows
- ▶ No additional bucket space allocated for future growth
- ▶ Negligible directory space
- ▶ Minor reorganization for splitting (redistribution occurs in the overflowed bucket only)

## ➤ Disadvantage

- ▶ Two block accesses: for directory, and for the bucket
- ▶ One block access in static hashing.

# Linear Hashing

- Dynamic expansion/shrinking of buckets without a directory
- Maintain overflow chains for each bucket
- For every overflow, buckets are split in the linear order.
- The overflowed bucket will eventually be split by the linear order → **delayed split**.
- Any records hashed to **bucket  $k$**  based on  $h_i$  will hash to **bucket  $k$**  or **bucket  $k+M$**  based on  $h_{i+1}$ .  
(Ex)  $h_1(r) = r \bmod M$  /  $h_2(r) = r \bmod 2M$  /  $h_3(r) = r \bmod 4M$

# Operations

- When a collision occurs,
  - ▶ put the record into its overflow chain
  - ▶ split the bucket  $k$  (*starting from 0*) pointed by  $n$  into bucket  $k$  and bucket  $k+M$
  - ▶ redistribute the records into the split buckets( $k, k+M$ ) using another hash function with  $h_{i+1}$
- To retrieve a record with key  $K$ 
  - ▶ apply hash function  $h_i$
  - ▶  $h_i(K) < n$ , apply  $h_{i+1}(K)$
- When  $n = M$ 
  - ▶ replace hash function and initialize  $n$  to 0

# Search Algorithm for linear hashing

**if**  $n = 0$

**then**  $m \leftarrow h_j(k)$

**else begin**

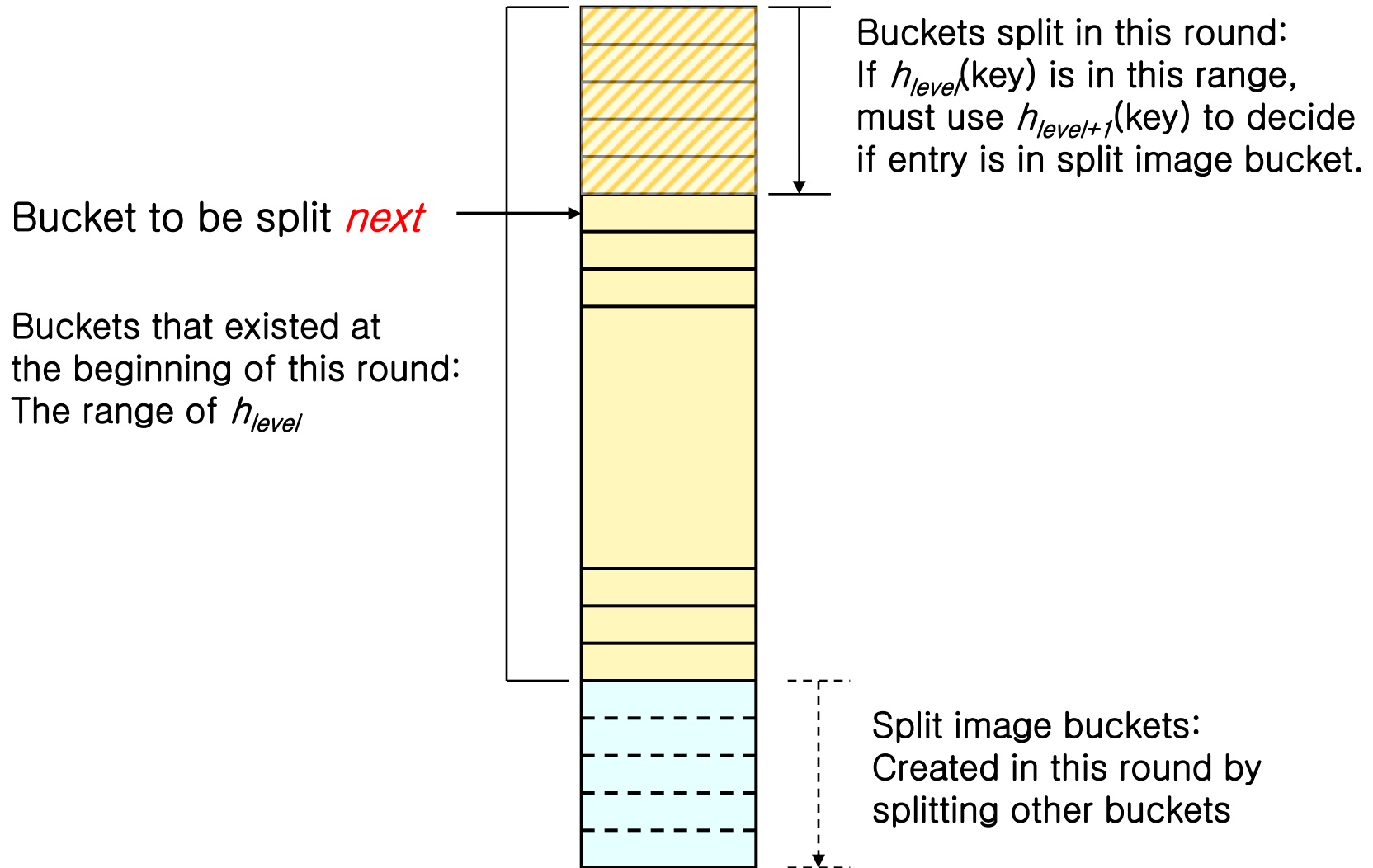
$m \leftarrow h_j(k)$

**if**  $m < n$  **then**  $m \leftarrow h_{j+1}(k)$

**end**

search the bucket whose hash value is  $m$   
( and its overflow if any)

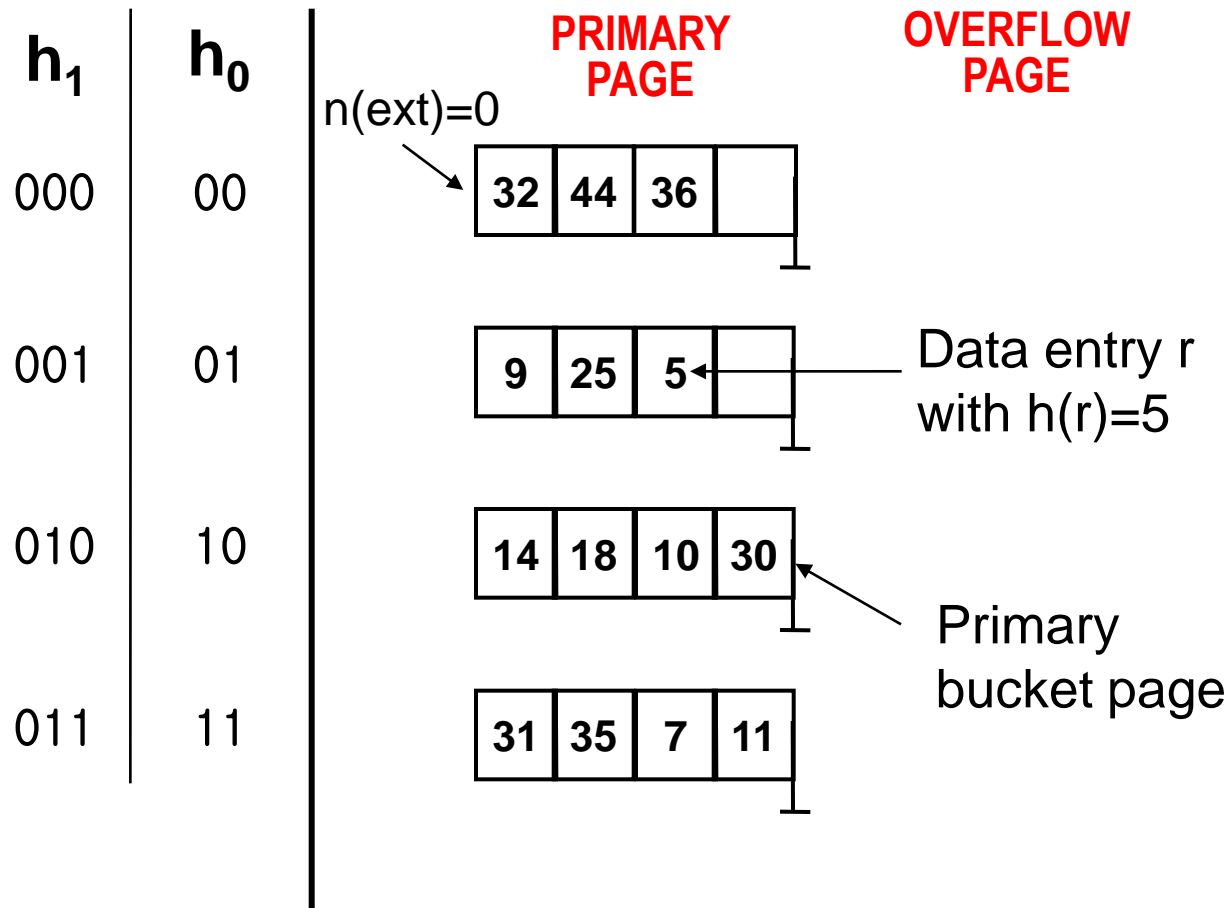
# Buckets during a Round





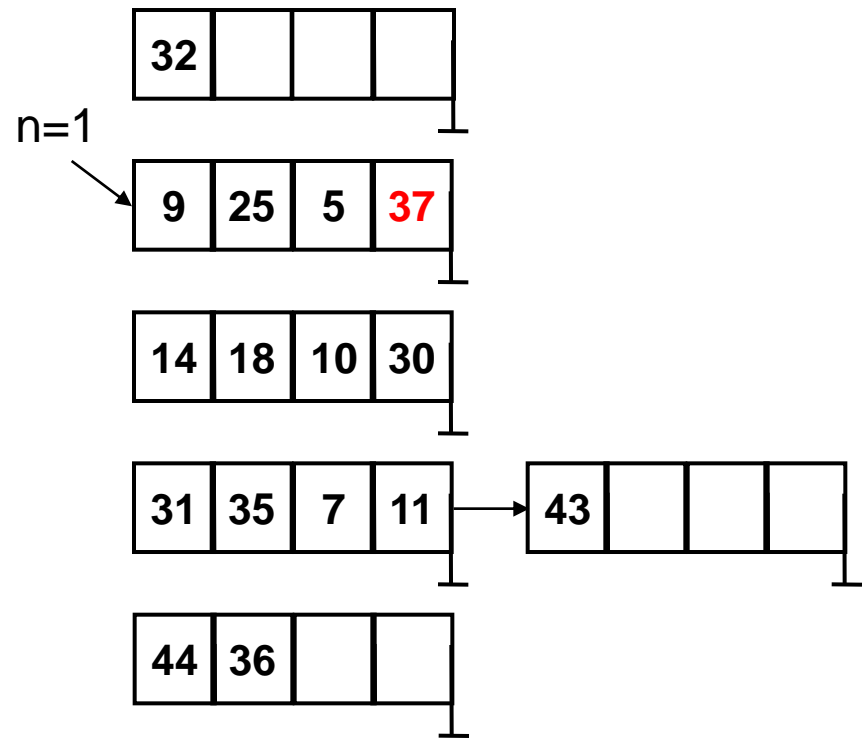
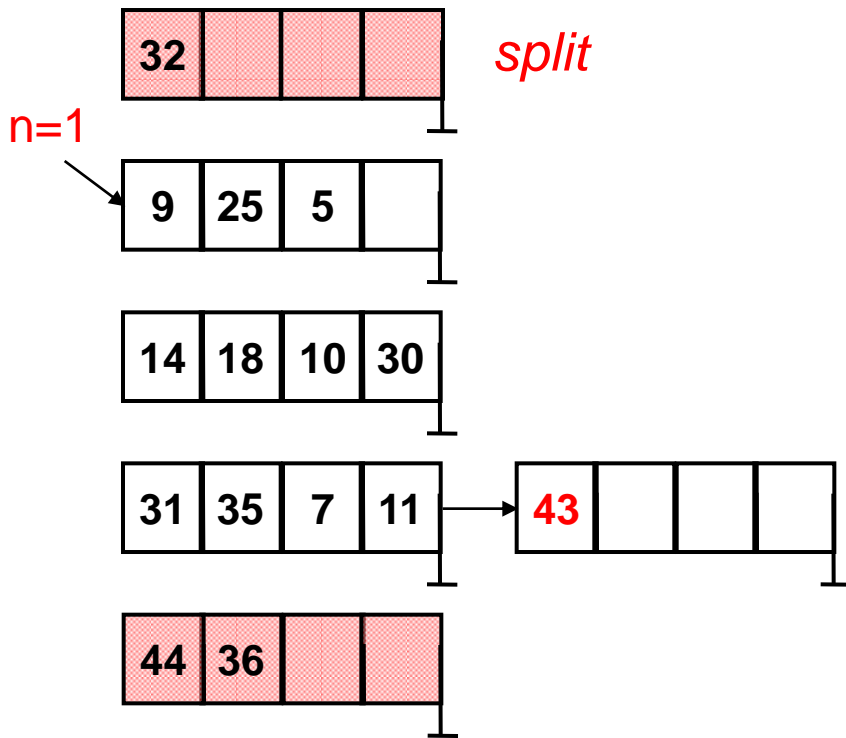
# Linear Hashing 예제(1)

Level=0, N=4



# Linear hashing 예제(2)

Level=0

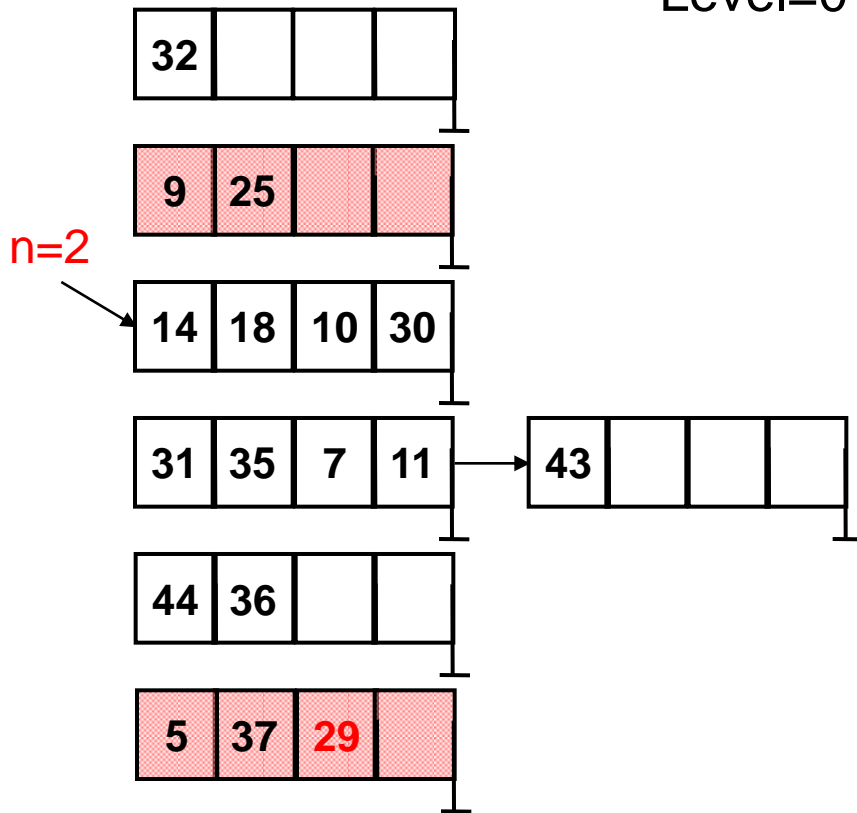


(a) Insert a record with 43

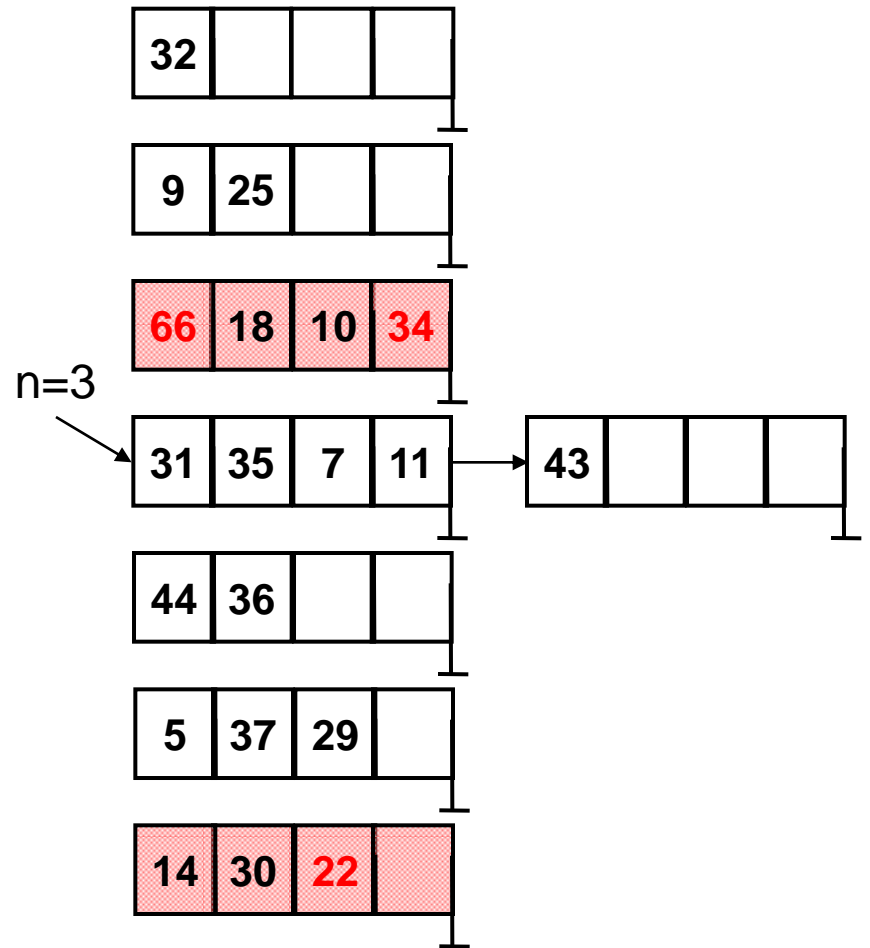
(b) Insert a record with 37

# Linear hashing 예제(3)

Level=0



(c) Insert a record with 29



(d) Insert a record with 22,66,34

# Linear hashing 예제(4)

