### SPEEDMAC: Speedy and Energy Efficient Data Delivery MAC Protocol for Real-Time Sensor Network Applications

**ICC 2010** 





### Motivation

### Sleep delay is the dominant factor of WSN packet latency

- A packet can traverse at most a single hop each cycle
  - Minimum packet latency = cycle time \*hops

### Most of WSN applications have real-time characteristics

- Disaster monitoring, real-time target tracking, intrusion detection, health, etc.
  - However, it is practically *impossible to obtain both low latency and low energy communication* at the same time

### Sleep delay exists for both synchronous & asynchronous MAC

- Synchronous scheduling (S-MAC, A-MAC)
  - A packet can traverse at most a single hop (or 2 with 'adaptive listening') each cycle since nodes beyond one-hop from the receiver cannot overhear the data.
- Asynchronous scheduling (B-MAC, Wise-MAC, XMAC)
  - A packet can traverse at most a single hop each cycle since a sender needs to send the preamble before starting the next-hop communication



# Motivation

### Synchronous skewed wakeup (DMAC) may be a solution!

- Schedule the wakeup time of each node in a pipelined fashion in the direction of packet movement so that
  - No sleep delay during the packet movement

### Issues with synchronous skewed wakeup

- May fail to deliver the message when multiple sensors compete for the message delivery
  - A single event is likely to be detected by nearby multiple sensors
  - Multiple events may occur simultaneously, which leads to collisions and contentions
- More idle listening
  - Since a node must wake up during the entire DATA transmission period instead of RTS period as in SMAC
- May not be practically possible to use such wakeup scheduling techniques for real applications unless these issues are completely resolved.

### Synchronous Skewed Wakeup





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### Synchronous Skewed Wakeup





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### **SPEED MAC Ideas**

- **Goal: Can we achieve both low-energy and low-latency at the same time?**
- **1. A collision signal to detect multi-source events & for fast event delivery** 
  - A special control packet called SIGNAL packet is used. It has different electrical characteristics from background noise

#### 2. Separate event report period from data delivery period

- Faster event report using a short control signal
- Lower energy consumption for idle period
  - To further reduce both the latency and the energy consumption

#### **3. Adaptive wakeup for multi-source events**

- Fast pipelined data delivery for a single-source event
- ► Full wakeup and CSMA-based data delivery for a multi-source event
  - Full duty-cycle operation for high-bandwidth transmission
  - Use RTS/CTS for busy periods



### Synchronous Skewed Wakeup





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### **Issues with Synchronous Skewed Wakeup**

#### Assumptions

- Stationary sensor nodes and stationary sinks
- Many to one communication pattern from multiple sources to the sinks

#### Issues

- Contention
  - Only a single source can transmit the data and other sources may have to wait
- Collision
  - When multiple nodes transmit at the same time, the packets will eventually collide in an upper layer and no packet can be transmitted
- Transmission error
  - When a transmission error occurs, the sender needs to wait for the next cycle

### For single-source event

- No contention, no collision, only need to consider error
- For multiple-source events
  - Need to consider contention, collision, and error



### **SPEED-MAC**

### Event announcement period: <u>Fast Event Announcement</u>

- In this period, nodes announce the presence of an event by sending a small control packet called a SIGNAL packet.
  - SIGNAL packet: consists of receiver address and collision bit
  - There is NO ACK packet for the signal packet.
- Collision detection for multi-source events
  - The *collision bit* tells that the event is a multi-source event.
  - Need to distinguish transmission errors from collision
- All the senders overhear the signal transmission from its parent
  - To distinguish a single source event from a multi-source event

### Data transmission period: <u>Adaptive Wakeup</u>

- In this period, nodes transfer messages by sending DATA packets
- ► For a single-source event, the period consists of DATA and ACK
  - Fixed scheduled data transmission for single-source events (not a CSMA)
- ► For a multi-source event, the period consists of RTS/CTS/DATA/ACK
  - Contention-based data transmission for multi-source events (CSMA/CA)

### **SPEED-MAC:** Single Source Event

#### No traffic

Nodes wakeup only during a signal rx slot.



- Single source traffic: single-packet data
  - Nodes wake up during signal rx/tx/rx slots and data slot





**PEED-MAC: Multi-Packet & Multi-Source Event** 

#### Single source traffic: multi-packet data

Nodes wake up during signal rx/tx/rx slots and multiple data slots

 Gepth: 0
 Gepth: 1

 Gepth: 2
 Gepth: 3

Data Packet

#### Multi-source traffic

► Nodes wake up during signal rx/tx/rx slots and several RTS/CTS/DATA/ACK slots

ACK packet





# **SPEED-MAC** with Multiple Sinks

### We can handle sink-to-sensor, sensor-to-sensor, and many sensors-to-many sinks scenarios





### **Collision/Error Differentiation**

#### Transmission error can occur due to two reasons

- Noise (Error)
  - Unwanted electrical signals interfering with the desired signal
  - The strength of the signal is irregular and variable
- Collision
  - Multiple simultaneous transmission collide at the receiver
  - The strength of the signal is regular and stronger
  - Can be differentiated at the physical layer by tracking RSSI
- In case of collision, the SIGNAL control packet is already destroyed.
  - COLLISION SIGNAL does not contain the receiver address anymore.
  - COLLISION SIGNAL packet is broadcast to the nodes in the upper layers
    - False-positive delivery: Nodes in the upper layers after the collision may unnecessarily wakeup

### **Collision/Error Differentiation**



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# **NS-2 Simulation Parameters**

- # of nodes: 400 grid nodes + 1 sink node
- Power
  - Tx : 30mW, Rx : 15mW, Idle : 15mW
- Bandwidth: 20Kbps
- Packet size
  - Data packet: 100B
  - Signal packet: 6B
  - Control packet: 10B
- Tx & Rx slot length
  - Data: 103ms, Signal: 22ms
- Simulation time: 10 min
- Total number of event: 20 events
- # of source nodes: 1, 2, 4, 8, 16 nodes
- Basic cycle time
  - SMAC: 1.44s
  - SPEED-MAC, D-MAC: 2.88s



# **Single Source – Latency**

#### **SMAC**

 SMAC suffers from the sleep delay and the additional buffering delay when the message generation interval is small.

#### SPEED-MAC vs. DMAC

- Due to the signaling wakeup period, SPEED-MAC's data latency is slightly higher than that of DMAC.
- Signal delivery latency of SPEED-MAC is almost close to the minimum delay achievable and is much smaller than DMAC's data delivery latency



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# **Single Source - Energy**

#### SMAC

As the packet generation interval decreases SMAC spends more energy in repeated wakeups and buffering.

#### SPEED-MAC vs. DMAC

- SPEED-MAC can achieve <u>an order of magnitude reduction</u> in the energy consumption compared to DMAC
  - By reducing the idle listening overhead and
  - By removing unnecessary wakeups during idle periods



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# **Multiple Sources - Latency**

#### **SMAC**

- Latency increases substantially as the number of source nodes increases. ►
  - This is due to the increased contention and buffering for multiple transactions.

#### **SPEED-MAC vs. DMAC**

- Constant and faster signal delivery latency even in multi-source events ►
- Noticeably higher data packet delay due to its adaptive wakeups and increased ► control packet (RTS and CTS) overhead for multi-source events.
  - For DMAC we use their assumption that an interference range of a node is twice







# **Multiple Sources - Energy**

#### SMAC

SMAC spends more energy due to its higher duty cycle operations

#### > SPEED-MAC vs. DMAC

Like the single-source case, SPEED-MAC can substantially reduce the energy consumption by reducing the idle listening and removing unnecessary wakeups.



### **MICA-2 Mote Implementation**

- Packet size: control packet: 10B, data packet: 100B
- Contention window: SYNC packet: 15 slots, Data packet: 31 slot

#### SINGLE SOURCE RESULTS

#### **MULTIPLE SOURCE RESULTS**



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